

THE POWER OF COMPUTATION

Lucy Martinez

Rutgers University

October 18, 2023

Advisor: Dr. Doron Zeilberger

About Me



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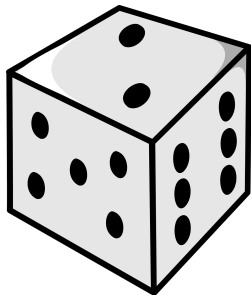
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GAME 1: DICE PROBLEM

Dice Rolls

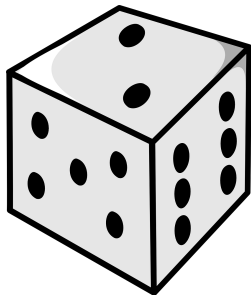
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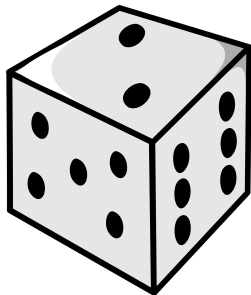
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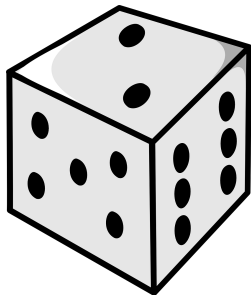
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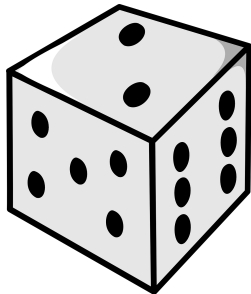


Suppose you have a fair standard die:

- Faces: $\{1, 2, 3, 4, 5, 6\}$
- Set $s = 0$. Roll the die and add the outcome to s
- Keep rolling the die, add the outcome to s , and stop when s is a prime number

Example

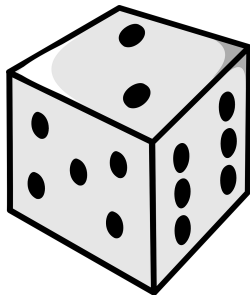
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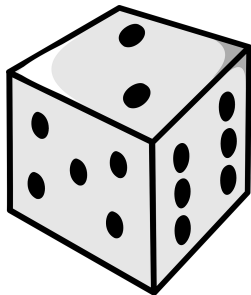
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- Say we get a 4, then $s = 4$



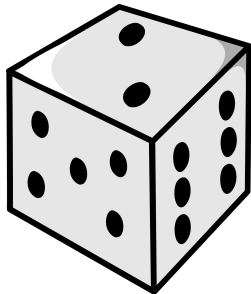
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- Say we get a 4, then $s = 4$
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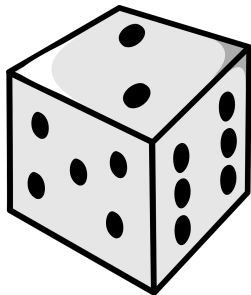
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- Roll again: say we get a 3, then $s = 10 + 3 = 13$ (prime!)

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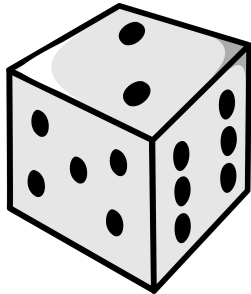
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Total rolls: 3

Example 2

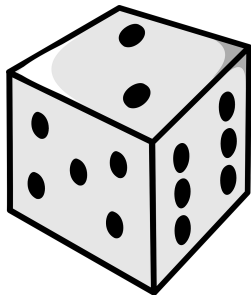
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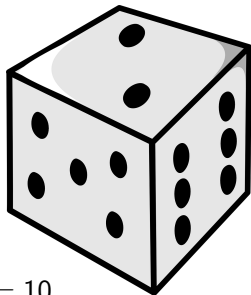
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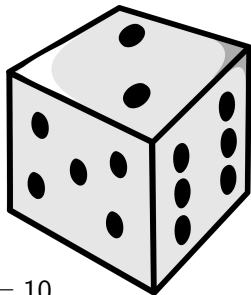
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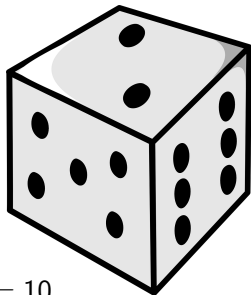
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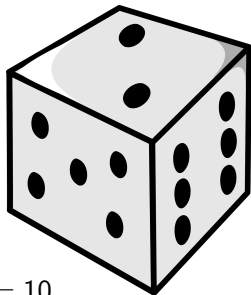
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- We would not be able to get a prime sum.

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- First roll: $\{2, 3, 5\}$, and with probability $1/2$ the game lasts one round

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With probability $1 - 1/2 - 2/9 = 5/18$, you need to continue

Estimate

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Alon-Malinovsky (2022)

The expectation of this random variable (up to an additive error of less than 10^{-4}) is 2.484

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- Turns out: $E_{1000} = 2.4284$ is a good approximation (Alon-Malinovsky)

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- What if instead of trying to hit a prime, you want to hit your favorite numbers? Say a product of two distinct primes, product of three distinct primes, perfect square (starting at a non-square), etc.

OUR APPROACH: SYMBOLIC COMPUTATION

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We used Maple to implement this function.

Symbolic Computation (continued)

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$$F_R(t, x) := F_{R-1}(t, x) + N_R(x)t^R.$$

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$$S_1(x) = P(x)S_0(x) - N_1(x) = \frac{1}{6}(x + x^4 + x^6)$$

$$\implies F_1(t, x) = \left(\frac{1}{6}(x^2 + x^3 + x^5)\right) t$$

First Two Rounds (continued)

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$$N_2(x) = \mathcal{P}(P(x)S_1(x))$$

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- If $s = 6$, we must roll: $\{1, 5\}$. Probability to get a prime sum is $1/6 \cdot 2/6 = 1/18$

The probability that the game lasts 2 rounds is

$$1/9 + 1/18 + 1/18 = \mathbf{2/9}$$

- First roll: $\{2, 3, 5\}$, and with probability $1/2$ the game lasts one round
- Second roll:

Possible outcomes in the first round: $\{1, 4, 6\}$

How can we get a prime sum if $s = 1$, $s = 4$ or $s = 6$?

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Note: The coefficient of $F_2(t, x)$ at t^2 was

$$\frac{1}{36}x^2 + \frac{1}{36}x^3 + \frac{1}{18}x^5 + \frac{1}{12}x^7 + \frac{1}{36}x^{11}$$

Non-rigorous Estimates - Results

Number of Faces	Property	Expected Duration
7	prime sum	2.1364...
12	prime sum	3.0814...
6	product of two distinct primes	3.7889...
6	product of three distinct primes	17.616887...
6	product of four distinct primes	112.907872...
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Note: To find the expected duration, we compute the partial derivative with respect to t of $F_R(t, x)$, evaluate at $t = x = 1$, and then divide by $F_R(1, 1)$.

GAME 2: ST. PETERSBURG PARADOX

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We would want to pay any amount A since $\infty - A = \infty$.

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In this case, if the gambler pays any amount A , then to ensure they do not lose money, $A < k + 1$.

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Example

Say $n = 6$ and $A = 0$. Then, $\text{StPetePT}(6,0)$ outputs

$[[2, 1/2], [4, 1/4], [8, 1/8], [16, 1/16], [32, 1/32], [32, 1/32]]$

Simulation (continued)

Next, we simulate the game in Maple:

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$M = [[2, 1/2], [4, 1/4], [8, 1/8], [16, 1/16], [32, 1/32], [32, 1/32]]$,
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Spoiler Alert: Using symbolic computation, the exact probability is $0.9088\dots$

Let $M = [[M_1, p_1], [M_2, p_2], \dots, [M_r, p_r]]$. Assume M_1, \dots, M_r are integers.

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we get

$$P_M(x) = \frac{1}{2}x^2 + \frac{1}{4}x^4 + \frac{1}{8}x^8 + \frac{1}{16}x^{16} + \frac{1}{16}x^{32}$$

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we get

$$P(x)^+ = \frac{1}{8} + \frac{1}{16} + \frac{1}{16} = 1/4$$

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Example

Playing the following

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for $n = 100$ times, we get the exact probability of 0.9088286275.

Advanced Computation

Essentially, we are interested in calculating

$$(P_M(x)^n)^+ = \sum_{j=1}^{\infty} \text{Coeff}_{x^j}(P_M(x)) = \cdots = \frac{1}{2\pi i} \int_{|x|=1} \frac{(P_M(x))^n}{x(x-1)} dx$$

for $n \in \mathbb{N}$, where $\text{Coeff}_{x^j}(P_M(x))$ is the coefficient of x^j in $P_M(x)$.

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Using this *theorem*, we can get a good approximation for sufficiently large n .

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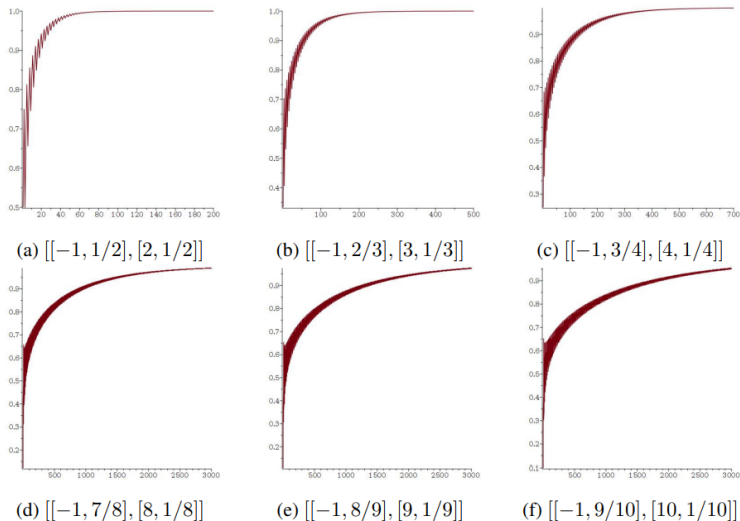


Figure 1. The risk-averseness graphs for the corresponding gambles.

- Dice Game and St. Petersburg Paradox

Conclusion

- Dice Game and St. Petersburg Paradox
- Simulation, and Symbolic Computation

THANK YOU!

Resources I



Lucy Martinez and Doron Zeilberger.

How many dice rolls would it take to reach your favorite kind of number?

To appear in Maple Transactions, 2023.



Lucy Martinez and Doron Zeilberger.

A guide to the risk-averse gambler and resolving the st. petersburg paradox once and for all.



Noga Alon and Yaakov Malinovsky.

Hitting a prime in 2.43 dice rolls (on average).

The American Statistician, 2023.

Alon-Malinovsky

For $k \leq n \leq 6k$ where n is non-prime, define $p(n, k)$ to be the probability that after k rolls, the running sum is n . Then,

Numerical Dynamic Programming

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$$p(k, n) = \frac{1}{6} \sum_i p(k-1, n-i)$$

where $i \in \{1, 2, \dots, 6\}$ such that $n-i$ is non-prime.

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Observe that $p(1, 1) = p(1, 4) = p(1, 6) = 1/6$.

Building Intuition

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$$E(\tau) = \sum_{k \geq 1} p(k) = p(1) + p(2) + \cdots$$


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It turns out,

$$p(k+1) := \sum_{n: k \leq n \leq 6k} p(k, n), \text{ non-prime } n$$

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Why? For non-prime n :

$$\begin{aligned} p(3) &= \sum_{3 \leq n \leq 18} p(2, n) \\ &= p(2, 4) + p(2, 6) + \dots + p(2, 18) \end{aligned}$$